

Tutorial Sheet - More on Branching and  
Kernelization  
NIT Calicut, July 2019

1. Show that any graph on  $n$  vertices with minimum degree 3 has a cycle of length at most  $2 \log n + 1$ . Can you also find such a cycle in  $O(m+n)$  time?
2. Show that  $(\log n)^k \leq (k \log k)^k + n$  for all positive integers  $n$  and  $k \leq n$ .
3. A tournament is a directed graph with a directed edge between every pair of vertices.
  - (a) Show that a tournament has a directed cycle if and only if it has a triangle.
  - (b) Consider the directed feedback vertex set problem where we are given a directed graph  $G$ , and an integer  $k$  and the goal is to determine whether there are  $k$  vertices whose removal results in a directed acyclic graph. Suppose the given directed graph is a tournament. Give a  $O^*(3^k)$  algorithm for the problem.
4.
  - (a) In the graph coloring problem, given a graph, and an integer  $k$  we want to determine if its vertices can be colored with at most  $k$  colors so that the end points of every edge gets different colors. What would be a brute force algorithm for the problem, and what is its running time? Is it possible to get a  $f(k)n^c$  time for the problem, where  $c$  is independent of  $k$ ?
  - (b) In the DUAL COLORING problem, we are given a graph  $G$  on  $n$  vertices and a positive integer  $k$ , and the objective is to test whether there exists a coloring of its vertices with at most  $n - k$  colors such that the end points of an edge gets different colors. Obtain a kernel with  $\mathcal{O}(k)$  vertices of this problem using crown decomposition (and crown lemma). Hint: Consider a crown decomposition in the complement of  $G$  and then infer what this decomposition looks like in the original graph, to design a reduction rule that deletes the head and crown of the crown decomposition.

5. In this problem, we will try to obtain a kernel for the  $d$ -hitting set problem. A sunflower with  $k$  petals and a core  $Y$  is a collection of sets  $S_1, S_2, \dots, S_k$  such that  $S_i \cap S_j = Y$  for all  $i \neq j$ ; the sets  $S_i \setminus Y$  are petals and we require none of them to be empty. Note that a family of pairwise disjoint sets is a sunflower (with an empty core).
- (a) Let  $A$  be a family of sets (without duplicates) over a universe  $U$ , such that each set in  $A$  has cardinality exactly  $d$ . Then show that if  $|A| > d!(k-1)^d$ , then  $A$  contains a sunflower with  $k$  petals and such a sunflower can be computed in time polynomial in  $|A|, |U|$  and  $k$ . (Try induction on  $d$ .)
  - (b) Can you come up with a reduction rule if there is a sunflower with  $k+1$  petals in the given  $d$ -hitting set instance?
  - (c) Obtain a kernel for the  $d$ -hitting set problem with  $O(d!k^d)$  sets.